Spectral and combinatorial properties of lexicographic powers of graphs

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Joint work with Nair Abreu², Paula Carvalho¹, Cybele Vinagre³

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Outline





- The spectra
- The Laplacian spectra

3 Some combinatorial properties of the lexicographic powers of graphs

4 References

Notation

- G = (V(G), E(G)) is a simple graph of order *n* and size *m*, vertex set V(G) and edge set E(G).
- $A_G = (a_{ij})$ is the adjacency matrix of G, that is, is the $n \times n$ matrix with $a_{ij} = \begin{cases} 1 & \text{if } ij \in E(G) \\ 0 & otherwise \end{cases}$;
- D_G is the diagonal matrix whose diagonal elements are the degrees d_1, \ldots, d_n of the vertices of G;
- $L_G = D_G A_G$ is the Laplacian matrix of G.

- The eigenvalues of A_G and L_G are indexed in nonincreasing order, i.e.,
 - $\lambda_1(G) \geq \lambda_2(G) \geq \cdots \geq \lambda_n(G);$
 - $\mu_1(G) \ge \mu_2(G) \ge \cdots \ge \mu_n(G) = 0.$

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• The all one vector j is the eigenvector of L_G associated to $\mu_n(G)$ and the multiplicity of $\mu_n(G)$ is equal to the number of components of G.

Definition

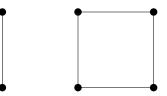
- its vertex set is the cartesian product $V(H) \times V(G)$
- and $(x_1, y_1) \sim (x_2, y_2)$ whenever $x_1 \sim x_2$ or $(x_1 = x_2 \text{ and } y_1 \sim y_2)$.

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The *lexicographic product* of two graphs H and G is the graph $H \cdot G$ (also called the graph *composition* and denoted H[G]) where

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This operation was introduced by Harary(1959) and Sabidussi (1959).

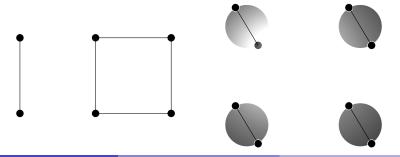


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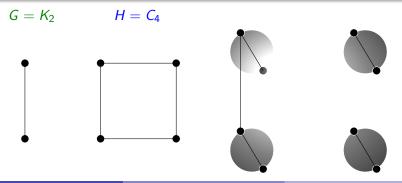
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The lexicographic product is associative but it is not commutative.



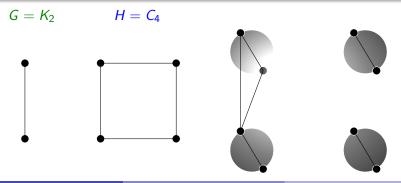
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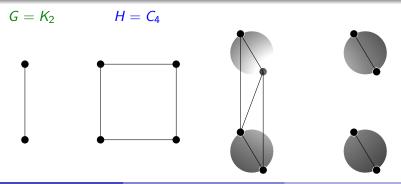
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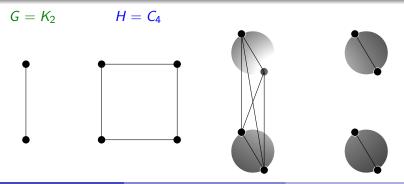


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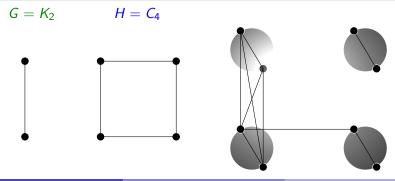
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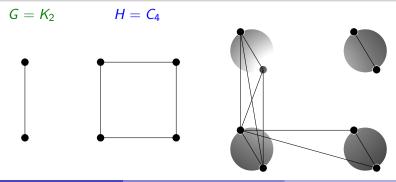
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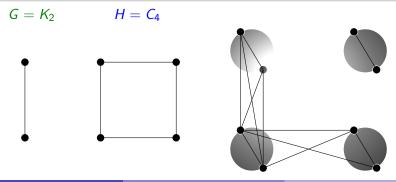
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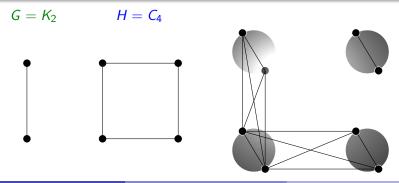
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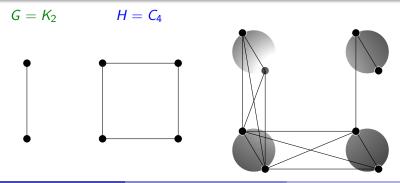
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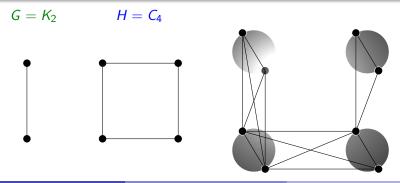
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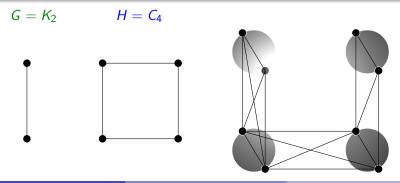
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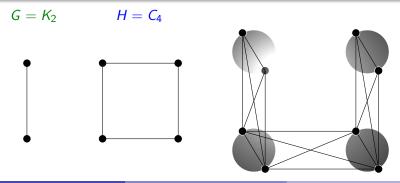
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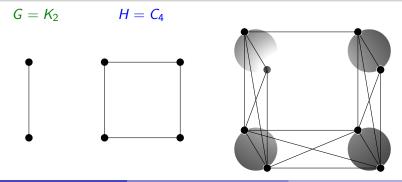
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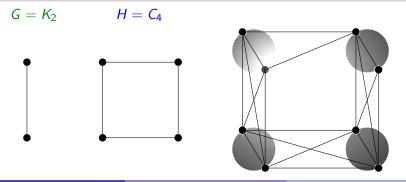
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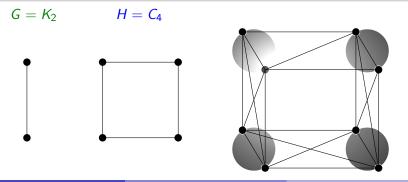
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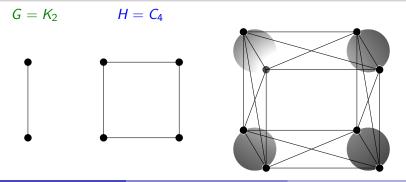
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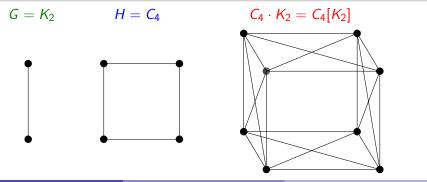
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Theorem

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$$\sigma_{A}(H[G]) = n(\sigma_{A}(G) \setminus \{p\}) \cup \sigma(mA_{H} + pI),$$

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 $\sigma_A(K_2) \setminus \{1\} = \{-1\}$ with multiplicity n = 4

$$\sigma_A(K_2) = \{1, -1\}$$

$$\sigma_A(C_4) = \{2, 0, 0, -2\}$$

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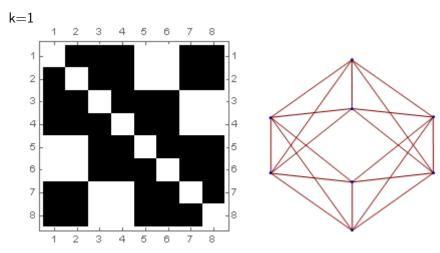
$$\begin{aligned} \sigma_A(K_2) &= \{1, -1\} & \sigma_A(K_2) \setminus \{1\} = \{-1\} \\ \sigma_A(C_4) &= \{2, 0, 0, -2\} & \text{with multiplicity } n = 4 \\ m &= 2 & \text{So} & 2 \times 2 + 1 = 5 \\ n &= 4 & 2 \times 0 + 1 = 1 \text{ (twice)} \\ p &= 1 & 2 \times (-2) + 1 = -3 \\ \sigma_A(C_4[K_2]) &= \{5, 1, 1, -3, -1, -1, -1, -1\} \end{aligned}$$

Arbitrary number of iterations of the lexicographic product

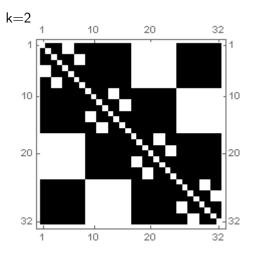
$$\begin{array}{rcl} H^0[G] &=& G, \\ H^1[G] &=& H[G] \\ H^2[G] &=& H[H[G]], \\ &\vdots \\ H^k[G] &=& H[H^{k-1}[G]] \end{array}$$

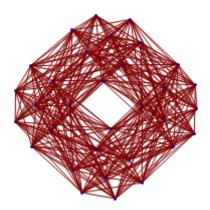
(we assume that the graph H is connected, and then, for $k \ge 1$, $H^k[G]$ is connected).

Example: Adjacency Matrix and graph for $C_4^k[K_2]$

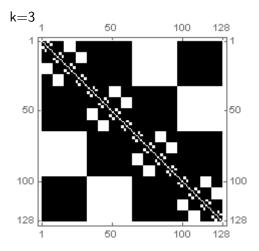


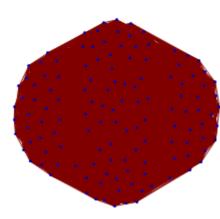
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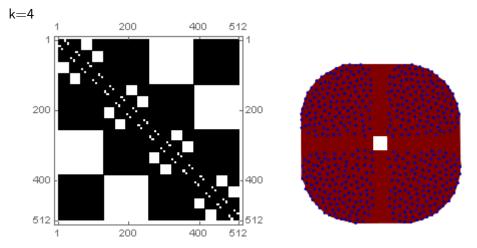


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Example: Adjacency Matrix and graph for $C_{4}^{k}[K_{2}]$



Iterated lexicographic products of graphs

Theorem

Let *G* be a connected *p*-regular graph of order *m* with *s* distinct eigenvalues such that $\sigma_A(G) = \{p, \lambda_2^{\lceil g_2 \rceil}(G), \ldots, \lambda_s^{\lceil g_s \rceil}(G)\}$ and let *H* be a connected *q*-regular graph of order *n* with *t* distinct eigenvalues such that $\sigma_A(H) = \{q, \lambda_2^{\lceil h_2 \rceil}(H), \ldots, \lambda_t^{\lceil h_t \rceil}(H)\}$, then $H^k[G]$ is a r_k -regular graph of order ν_k , such that

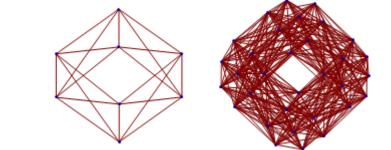
$$\mathbf{r}_{\mathbf{k}} = \mathbf{m}q\frac{n^{k}-1}{n-1} + \mathbf{p}, \qquad \mathbf{\nu}_{\mathbf{k}} = \mathbf{m}n^{k},$$

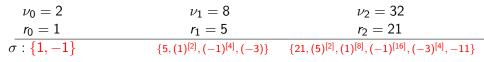
$$\sigma_{\mathcal{A}}(\mathcal{H}^{k}[G]) = \left\{\lambda_{2}^{[n^{k}g_{2}]}(G), \ldots, \lambda_{s}^{[n^{k}g_{s}]}(G)\right\} \cup \left\{mq\frac{n^{k}-1}{n-1}+p\right\} \cup \Lambda_{k}$$

where, for $k \ge 1$, $\Lambda_{k} = \bigcup_{i=0}^{k-1} \Big\{ (mn^{i}\lambda_{2}(H) + r_{i})^{[n^{k-1-i}h_{2}]}, \dots, (mn^{i}\lambda_{t}(H) + r_{i})^{[n^{k-1-i}h_{t}]} \Big\}.$

Example: the spectrum of $H^k[G] = C_4^k[K_2]$.

Notice that $H = C_4$ and $G = K_2(m = 2, n = 4, p = 1 \text{ and } q = 2)$ k = 0 k = 1 k = 2





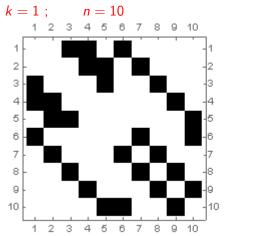
Spectra of powers of lexicographic products of graphs

Corollary

Let H be a connected q-regular graph of order n with t distinct eigenvalues such that $\sigma_A(H) = \{q, \lambda_2^{[h_2]}(H), \ldots, \lambda_t^{[h_t]}(H)\}$. Then H^k is a r_k -regular graph of order ν_k , such that

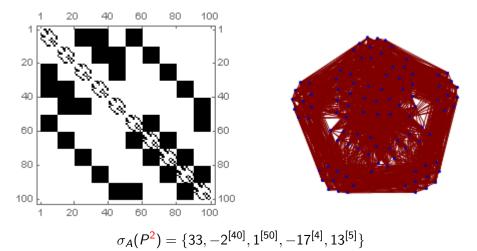
for $k \geq 1$.

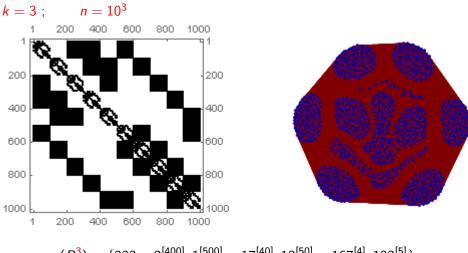
The number of distinct eigenvalues of H^k is not greater that (t-1)k+1.





k = 2; $n = 10^2$





 $\sigma_{\mathcal{A}}(\mathcal{P}^{\mathbf{3}}) = \{333, -2^{[400]}, 1^{[500]}, -17^{[40]}, 13^{[50]}, -167^{[4]}, 133^{[5]}\}$

And, ... for k = 10; $n = 10^{10}$

$\sigma_A(P^{10})$	=	$\Big\{ 3333333333, -2^{[4000000000]}, 1^{[5000000000]},$
		$-17^{[400000000]}, 13^{[500000000]},$
		$-167^{[40000000]}, 133^{[50000000]},$
		$-1667^{[4000000]}, 1333^{[5000000]},$
		$-16667^{[400000]}, 13333^{[500000]},$
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		$-16666667^{[400]}, 13333333^{[500]},$
		$-166666667^{[40]}, 133333333^{[50]},$
		$-1666666667^{[4]}, 1333333333^{[5]} \Big\}$

k	Spectrum of P ^k		
k = 1	$3, 1^{[5]}, -2^{[4]}$		
<i>k</i> = 2	33, 13[5], 1[50], -2[40], -17[4]		
<i>k</i> = 3	$333, 133^{[5]}, 13^{[50]}, 1^{[500]}, -2^{[400]}, -17^{[40]}, -167^{[4]}$		
k = 100	$3 \times \sum_{i=0}^{99} 10^{i}, \qquad 1^{[5 \times 10^{99}]}, \ -2^{[4 \times 10^{99}]},$ $\left(10^{m} + 3 \sum_{i=0}^{m-1} 10^{i} \right)^{[5 \times 10^{99-m}]}, \ m = 1, \dots, 99,$ $- \left(7 + 10^{m} + 6 \sum_{i=1}^{m-1} 10^{i} \right)^{[4 \times 10^{99-m}]}, \ m = 1, \dots, 99.$		

Notice that the graph P^k has 10^k vertices, in particular P^{100} has the googol number of vertices 10^{100} . All the computations were done by Mathematica and lasted just a few seconds.

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The Laplacian spectra of lexicographic compositions

From [C., Freitas, Martins, Robbiano, 2013] we can also deduce the following result.

Theorem

Let G be a graph of order m and H a graph of order n, then

$$\sigma_L(H[G]) = \left(\bigcup_{j=1}^n (md_H(j) + (\sigma_L(G) \setminus \{0\}))\right) \cup m\sigma_L(H).$$

The Laplacian spectra of lexicographic compositions

Theorem

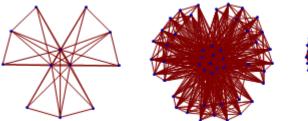
Let G be a connected graph of order m and H be a connected graph of order n. Then $H^k[G]$ is a graph of order $\nu_k = mn^k$, and

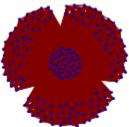
$$\sigma_L(\boldsymbol{H^k[G]}) = \Omega_{\boldsymbol{G}}^k \cup \Omega_{\boldsymbol{H}}^k,$$

where

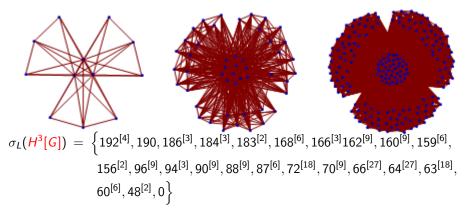
$$\Omega_{G}^{k} = \bigcup_{(j_{1},j_{2},\dots,j_{k})\in[n]^{k}} \left(\sigma_{L}(G)\setminus\{0\} + m\sum_{i=1}^{k} n^{i-1}d_{H}(j_{i})\right)$$
$$\Omega_{H}^{k} = \bigcup_{i=2}^{k} \left(\bigcup_{(j_{i},\dots,j_{k})\in[n]^{k-i+1}} \left(mn^{i-2}\sigma_{L}(H)\setminus\{0\} + m\sum_{r=i}^{k} n^{r-1}d_{H}(j_{r})\right)\right) \bigcup mn^{k-1}\sigma_{L}(H)$$
for $k \geq 1$.

Example: $H^k[G]$ with $H = K_{1,3}$, $G = P_3$ $\sigma_L(K_{1,3}) = \{4, 1, 1, 0\};$ $\sigma_L(P_3) = \{3, 1, 0\};$ n = 4; m = 3 k = 1 k = 2 k = 3 $\nu_1 = 12$ $\nu_2 = 48$ $\nu_3 = 192$





Example: $H^k[G]$ with $H = K_{1,3}$, $G = P_3$ $\sigma_L(K_{1,3}) = \{4, 1, 1, 0\};$ $\sigma_L(P_3) = \{3, 1, 0\};$ n = 4; m = 3 k = 1 k = 2 k = 3 $\nu_1 = 12$ $\nu_2 = 48$ $\nu_3 = 192$



Lexicographic powers of graphs

Now, let us consider the case G = H.

Corollary

Let *H* be a connected graph of order *n*. Then H^k is a graph of order $\nu_k = n^k$, such that

$$\sigma_{L}(\boldsymbol{H}^{k}) = \bigcup_{i=1}^{k-1} \left(\bigcup_{(j_{i},\dots,j_{k-1})\in[n]^{k-i}} \left(n^{i-1}\sigma_{L}(\boldsymbol{H})\setminus\{0\} + \sum_{r=i}^{k-1} n^{r}d_{H}(j_{r}) \right) \right) \bigcup n^{k-1}\sigma_{L}(\boldsymbol{H})$$

for $k \geq 1$.

Lexicographic powers of graphs

Example: $H = K_{1,3}$: n = 4, $\sigma_L(K_{1,3}) = \{4, 1, 1, 0\}$.

 $H^3 = K_{1,3}^3$ is a graph of order

$$v_3 = 4^3 = 64$$



with

$$\sigma_{L}(\mathcal{K}_{1,3}^{3}) = \bigcup_{i=1}^{2} \left(\bigcup_{(j_{i},j_{2})\in[4]^{3-i}} \left(4^{i-1}\sigma_{L}(\mathcal{K}_{1,3}) \setminus \{0\} + \sum_{r=i}^{2} 4^{r}d_{H}(j_{r}) \right) \right) \bigcup 4^{2}\sigma_{L}(\mathcal{K}_{1,3})$$

which gives $\sigma_L(K_{1,3}^3) =$

 $\left\{ 64^{[3]}, 61^{[2]}, 56^{[2]}, 53^{[6]}, 52^{[2]}, 32^{[6]}, 29^{[6]}, 24^{[9]}, 21^{[18]}, 20^{[6]}, 16^{[2]}, 0 \right\}$

Laplacian index

$$\mu_1(H^k) = n^{k-1}\mu_1(H).$$

• Algebraic connectivity

$$\mu_{n^{k}-1}(H^{k}) = n^{k-1}\mu_{n-1}(H).$$

• Minimum and maximum degree

$$\delta(H^k) = \delta(H) \frac{n^k - 1}{n - 1}$$
 and $\Delta(H^k) = \Delta(H) \frac{n^k - 1}{n - 1}$

Stability number

As $\alpha(H[G]) = \alpha(H) \alpha(G)$ [Geller, 1975], where G is an arbitrary graph

$$\alpha(H^k) = \alpha(H)^k.$$

Furthermore, from the spectral upper bound [(Godsil (2008) and Lu (2007)] for an arbitrary graph G

$$\alpha(G) \leq n \; \frac{\mu_1(G) - \delta(G)}{\mu_1(G)},$$

we obtain,

$$\begin{aligned} \alpha(H^k) &\leq n^k \frac{\mu_1(H^k) - \delta(H^k)}{\mu_1(H^k)} \\ &= n^k \frac{\mu_1(H) - \delta(H)}{\mu_1(H)}. \end{aligned}$$

• Vertex connectivity

As the lexicographic product H[G] is connected if and only if H is a connected graph [Harary and Wilcox, 1967] if both G and H are not complete [Geller and Stahl, 1975]

$$\upsilon(H[G]) = m \,\upsilon(H),$$

where v(H) denotes the vertex connectivity of H. So,

$$v(H^k) = n^{k-1}v(H).$$

Furthermore, we may conclude that when H is connected not complete (and then H^k is also connected not complete),

$$n^{k-1}\mu_{n-1}(H) \leq \upsilon(H^k) \leq \delta(H) \frac{n^k-1}{n-1}.$$

• Chromatic number

It is well known the following lower bound due to Hoffman

$$\chi(G) \geq 1-\frac{\lambda_1(G)}{\lambda_n(G)}$$

As direct consequence, if a graph H is q-regular of order n,

$$\chi(H^k) \geq 1 - \frac{r_k}{\lambda_{n^k}(H^k)}$$

= $1 - \frac{n^k - 1}{n^{k-1}\left((n-1)\frac{\lambda_n(H)}{q} + 1\right) - 1}$

Proposition

Let H be a connected not complete graph and let G be an arbitrary graph of order m. For every $k \in \mathbb{N}$

$$diam(H^{k+1}) = diam(H^k[G]) = diam(H).$$

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THANK YOU!



Ten years after, congratulations again Professor Dragos Cvetković

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