

Global Optimality Conditions and Methods for Nonconvex Mathematical Programming

Alexander S. Strekalovsky

Matrosov Institute for System Dynamics and Control Theory
of Siberian Branch
of Russian Academy of Sciences

e-mail: strelkal@icc.ru

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Definition

Let set $X \subset \mathbb{R}^n$ be convex. Then $f: X \rightarrow \mathbb{R}$ is a d.c. function on X , when

$$\exists g, h \in \text{Conv}(X): \quad f(x) = g(x) - h(x) \quad x \in X.$$

$DC(X)$ is a linear space of d.c. functions on X .

$\text{Conv}(X)$ is a convex cone of convex functions on X .

1) $DC(X) = \text{lin}(\text{Conv}(X))$.

$$\left\{ \begin{array}{l} \mathcal{K}(X) \text{ is a convex cone of functions} \\ \text{Conv}(X) \subset \mathcal{K}(X) \subset DC(X), \\ \mathcal{K}(X) - \mathcal{K}(X) = DC(X). \end{array} \right.$$

2) $C^2(\Omega) \subset DC(\Omega)$, where Ω is a open convex set.

3) $\text{cl}(DC(X)) = C(X)$, if X is a convex compact.

4) $DC(X)$ is a closed w.r.t. the following operations:

$$\sum_{i=1}^m \lambda_i f_i(x); \quad \max_i f_i(x); \quad \min_i f_i(x); \quad |f(x)|; \quad \prod_{i=1}^m f_i(x);$$

$$f^+(x) = \max \{0, f(x)\}; \quad f^-(x) = \min \{0, f(x)\}.$$

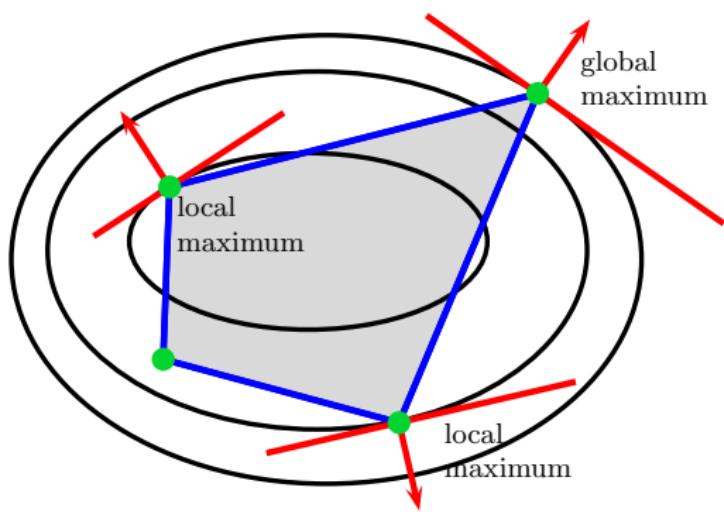


$$\begin{cases} f_0(x) = g_0(x) - h_0(x) \downarrow \min, & x \in S, \\ f_i(x) = g_i(x) - h_i(x) \leq 0, & i \in I = \{1, \dots, m\}. \end{cases} \quad (\mathcal{P})$$

$g_i, h_i \in \text{CONVEX}(\mathbb{R}^n), \quad i \in \{0, 1, \dots, m\},$

$S \subset \mathbb{R}^n$ is a closed convex set.

- ① Linearization of basic nonconvexities of the problem under scrutiny and, consequently, reduction of the problem to a family of (partially) linearized problems.
- ② Application of convex optimization methods for solving linearized problems and, as a consequence, “within” special local search methods.
- ③ Construction of “good” approximations (resolving sets) of level surfaces/epigraph boundaries of convex functions.



Practical Rules for Solving Nonconvex Optimization Problems

- ① Never apply convex optimization methods DIRECTLY.
- ② Exact classification of the problem under scrutiny.
- ③ Application of special (for the class of problem under scrutiny) local search (LS) methods, or (problem) specific methods.
- ④ Application of global search strategies specialized for the class of nonconvex problems.
- ⑤ Construction of pertinent approximations of level surfaces with the aid of the experience obtained during solving similar nonconvex problems.
- ⑥ Application of convex optimization methods for solving linearized problems and within the framework of special LS methods.

- ① Special local search methods. Convergence theorems.
- ② Global Optimality Conditions (GOC).
- ③ Global search algorithms.
- ④ Convergence theorems of global search algorithms.
- ⑤ Testing.

A.S. Strekalovsky, *Elements of Nonconvex Optimization* (Nauka, Novosibirsk, 2003) [in Russian].

Example 1. Incorrect application of Newton's method

Problem Formulation

$$f(x) \triangleq \ln(1 + x_1) + \ln(1 + x_2) \downarrow \min, \quad x \in \Pi \subset \mathbb{R}^2,$$
$$\Pi = \{ (x_1, x_2) \mid -\frac{1}{2} \leq x_i \leq 3 \}.$$

It's obvious that the point $z = (-\frac{1}{2}, -\frac{1}{2})^T$ is the solution to this problem.

Consider the point $x^k = (0, 0)^T$,

$$\nabla f(x) = \left(\frac{1}{1+x_1}, \frac{1}{1+x_2} \right)^T, \quad \nabla f(x^k) = (1, 1)^T,$$

$$\nabla^2 f(x) = \begin{pmatrix} -\frac{1}{(1+x_1)^2} & 0 \\ 0 & -\frac{1}{(1+x_2)^2} \end{pmatrix}, \quad \nabla^2 f(x^k) = \begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}.$$

Auxiliary problem of Newton's method

$$\phi(d) = d_1 + d_2 - d_1^2 - d_2^2 \downarrow \min, \quad d \in D = (D - x^k).$$

has the solution $d = (3, 3)^T$, which is a direction to the worst feasible point $x = (3, 3)^T$.

Пусть $f, g, h \in CONVEX(\mathbb{R}^n)$.

D.C. minimization

$$g(x) - h(x) \downarrow \min, \quad x \in D. \quad (DC)$$

D.C. constrained problem

$$\left. \begin{array}{l} f(x) \downarrow \min, \quad x \in S, \\ g(x) - h(x) \leq 0. \end{array} \right\} \quad (DCC)$$

Note:

- i) If $g \equiv 0$ in (DC) , then (DC) is convex maximization problem;
- ii) If $g \equiv 0$ in (DCC) , then (DCC) is reverse-convex problem.

Convex Maximization Problem

Problem formulation

$$h(x) \uparrow \max, \quad x \in D. \quad (P)$$

Theorem 1. (Global optimality Conditions)

Let $z \in D$ be a global maximum to problem (P) ($z \in Sol(P)$).

Then,

$$\left. \begin{array}{l} \forall y : h(y) = h(z), \quad \forall x \in D \\ \langle \nabla h(y), x - y \rangle \leq 0. \end{array} \right\} \quad (1)$$

If, in addition, the following regularity condition holds

$$(H) : \quad \exists v \in I\!\!R^n : -\infty < h(v) < h(z). \quad (2)$$

then condition (1) turns out to be sufficient for the point z being a global solution to Problem () .

1) If $z \in Sol(P)$, then by the convexity $h(\cdot)$ the following condition holds:

$$\forall y : h(y) = h(z), \quad \forall x \in D : \quad 0 \geq h(x) - h(z) = h(x) - h(y) \geq \langle \nabla h(y), x - y \rangle.$$

2) Now suppose there exists a vector $u \in D : h(u) > h(z)$. Consider the Lebesgue set of function $h(\cdot)$

$$S(h, z) = \{x \in I\!R^n \mid h(x) \leq h(z)\}.$$

The set $S(h, z)$ is a convex and closed. Besides, $v \in int S(h, z)$. Then there is a projection the point u on set $S(h, z)$, i.e.

$$\exists y \in S(h, z) : \quad h(y) = h(z),$$

$$\frac{1}{2} \|y - u\|^2 = \inf_x \left\{ \frac{1}{2} \|x - u\|^2 : \quad x \in S(h, z) \right\}.$$

Since $u \notin S(h, z)$, then

$$\|y - u\| > 0. \tag{3}$$

By KKT-theorem the projection y satisfying the following condition:

$$\left. \begin{array}{l} \lambda_0(y - u) + \lambda \nabla h(y) = 0, \quad \lambda_0, \lambda \geq 0, \\ \lambda_0 + \lambda > 0, \quad \lambda(h(y) - h(z)) = 0. \end{array} \right\} \quad (4)$$

If $\lambda_0 = 0$, then $\lambda > 0$ and $\nabla h(y) = 0$. By the convexity of $h(\cdot)$, it means that $y \in \text{Argmin}(h, I\!\!R^n)$. But this contradicts assumption (2), because $h(y) = h(z)$.

Now if $\lambda = 0$, then $\lambda_0 > 0$, and $y = u$ follows from (4). It is impossible in view of (3). Further, dividing the (4) by λ , we get

$$\lambda_0(y - u) + \nabla h(y) = 0, \quad \lambda_0 > 0.$$

Hence (3) implies the following inequality

$$\langle \nabla h(y), u - y \rangle = \lambda_0 \|y - u\|^2 > 0,$$

which contradicts (1).

Relations with Classical Optimality Conditions

- a) From the Global Optimality Condition (1) with $y = z$ it follows necessary condition of local optimality for problem (P) :

$$\langle \nabla h(z), x - z \rangle \leq 0 \quad \forall x \in D, \quad (5)$$

- b) Let the feasible set D is defined by the differentiable functions, as follows,

$$D = \{x \in I\!\!R^n \mid g_i(x) \leq 0, i = 1, \dots, m\}. \quad (6)$$

Then, it is easy to show that the KKT-condition follows from (5), because (5) implies that z is a solution to problem:

$$\langle \nabla h(z), x \rangle \uparrow \max, \quad g_i(x) \leq 0, \quad i = 1, \dots, m.$$

$$y : \quad f(y) = f(z) \quad \exists(\lambda_0, \lambda_1, \dots, \lambda_m) \neq 0 :$$

$$\left. \begin{aligned} & -\lambda_0 \nabla h(y) + \sum_{i=1}^m \lambda_i \nabla g_i(x(y)) = 0, \\ & \lambda_i \geq 0, \quad \lambda_i g_i(x(y)) = 0, \quad i = 1, \dots, m. \end{aligned} \right\} \quad (7)$$

For all $y : h(y) = h(z)$ there exists Lagrange multipliers vector

$$\Lambda(y) = (\lambda_0 = \lambda_0(y), \quad \lambda_i = \lambda_i(y)).$$

Additionally, along with conditions (7) it is necessary to check the following inequality

$$\langle \nabla h(y), x(y) - y \rangle \leq 0, \quad i = 1, \dots, m,$$

where $x(y) \in D$ is a solution to linearized problem:

$$\langle \nabla h(y), x \rangle \uparrow \max, \quad g_i(x) \leq 0, \quad i = 1, \dots, m.$$



- c) It can be shown that the Rockafellar condition

$$\partial h(z) \subset N(z|D) \quad (8)$$

follows from GOC with $y = z$.

- d) The Global Optimality Conditions (1) has the algorithmic property.

Suppose that y , $h(y) = h(z)$ and $u \in D$ are such that $\langle \nabla h(y), u - y \rangle > 0$, then in view of the convexity of $h(\cdot)$ we get $h(u) > h(y) = h(z)$.

- e) Verification of the Global optimality Conditions is reduced to solving family of linearized problems (for all y : $h(y) = h(z)$):

$$\langle \nabla h(y), x \rangle \uparrow \max, \quad x \in D, \quad (PL(y))$$

and verification of inequality:

$$\langle \nabla h(y), x(y) - y \rangle \leq 0, \quad (9)$$

where $x(y) \in Sol(PL(y))$.

Example 2

Consider the following problem

$$h(x) = \begin{cases} (x_1^2 - x_2) \uparrow \max, & x \in D \subset \mathbb{R}^2, \\ D = D_1 \cup D_2, & \end{cases} \quad (10)$$

$$\begin{aligned} D_1 &= \{x = (x_1, x_2) \mid 0 \leq x_1 \leq 1, x_2 \geq 0\}, \\ D_2 &= \{x = (x_1, x_2) \mid -1 \leq x_i \leq 0, i = 1, 2\}. \end{aligned} \quad (11)$$

It is easy to see that the point $z = (1, 0)^T$ satisfies to the classical optimality condition ($\nabla h(z) = (2, -1)^T$)

$$\langle \nabla h(z), x - z \rangle \leq 0 \quad \forall x \in D,$$

in spite of nonconvexity of the feasible set D . On the other hand, for $y = (-1, 0)^T$, $h(y) = 1 = h(z)$, $y \in D_2 \subset D$, condition (1) is violated with $u = (-1, -0.5) \in D_2 \subset D$:

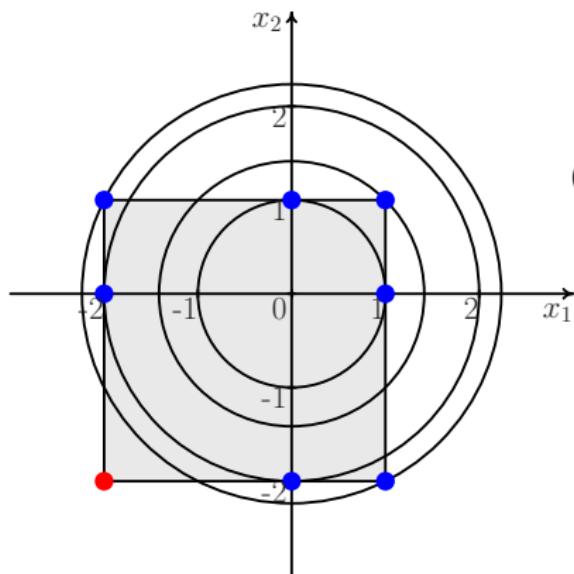
$$\langle \nabla h(y), u - y \rangle = \frac{1}{2} > 0.$$

Therefore, $z = (0, 1)^T$ is not a global solution to problem (10)–(11).

Example 3. Convex Maximization Problem

$$\|x\|^2 \uparrow \max_x, \quad x \in \Pi,$$

$$\Pi = \{ x \in \mathbb{R}^n \mid -2 \leq x_i \leq 1, \quad i = 1, 2, \dots, n \}.$$



2^n is a number of local maxima
 $(3^n - 1)$ is a number of points satisfying
Karush-Kuhn-Tucker conditions.

UNIQUE global maximum!!!



D.C. Minimization Problem

D.C. Minimization Problem. Problem Formulation

$$g(x) - h(x) \downarrow \min, \quad x \in D. \quad (\text{DC})$$

The basic element of the local search method is solving the (linearized at a current feasible point $x^s \in D$) convex problem

$$J_s(x) = g(x) - \langle h'_s, x \rangle \downarrow \min, \quad x \in D, \quad (\mathcal{PL}_s)$$

where $h'_s = h'(x^s) \in \partial h(x^s)$.

Next point $x^{s+1} \in D$ is constructed as an approximate solution of the problem (\mathcal{PL}_s) .

Theorem 2

Let $F = g - h$ be a bounded below function on D and function $h(\cdot)$ be convex on D .

Then the sequence $\{x^s\}$, generated by the rule

$$g(x^{s+1}) - \langle h'(x^s), x^{s+1} \rangle \leq \inf_x \{g(x) - \langle h'(x^s), x \rangle \mid x \in D\} + \delta_s, \quad (12)$$

satisfies the following condition:

$$\lim_{s \rightarrow \infty} [\mathcal{V}(PL_s) - J_s(x^s)] = 0. \quad (13)$$

At the same time any accumulation point x_* of sequence $\{x^s\}$ is a solution of the problem

$$J_*(x) = g(x) - \langle h'(x_*), x \rangle \downarrow \min, \quad x \in D, \quad (\mathcal{PL}(x_*))$$

where $h'(x_*) \in \partial h(x_*)$.

If $h(\cdot)$ is strongly convex, then $x^s \rightarrow x_* \in D$. Besides,

$$\|x^s - x^{s+1}\|^2 \leq \frac{2}{\mu} (F(x^s) - F(x^{s+1}) + \delta_s),$$

where $\mu > 0$ is a strong convexity constant of function $h(\cdot)$.



Problem Formulation

$$(\mathcal{P}): \quad f(x) \triangleq \ln(1 + x_1) + \ln(1 + x_2) \downarrow \min, \quad x \in \Pi \subset \mathbb{R}^2,$$

$$\Pi = \{ (x_1, x_2) \mid -\frac{1}{2} \leq x_i \leq 3 \}.$$

The solution of the problem (\mathcal{P}) is the point $z = (z_i = -\frac{1}{2}, \quad i = 1, 2, \dots, n)$.

$$f(x) = g(x) - h(x), \quad g \equiv 0.$$

$$\nabla f(x) = -\nabla h(x) = \left(\frac{1}{1+x_i}, \quad i = 1, 2, \dots, n \right), \quad x^0 \in \Pi.$$

$$(\mathcal{P}L_s): \quad \langle \nabla f(x^0), x \rangle = -\langle \nabla h(x^0), x \rangle \downarrow \min, \quad x \in \Pi.$$

$$\sum_{i=1}^n \left(\frac{1}{1+x_i^0} \cdot x_i \right) \downarrow \min, \quad -\frac{1}{2} \leq x_i \leq 3, \quad i = 1, \dots, n.$$

Therefore,

$$x^1 = z \triangleq \left(-\frac{1}{2}, -\frac{1}{2}, \dots, -\frac{1}{2} \right) \in Sol(\mathcal{P}).$$

General procedure of global search consists of the two parts:

a) local search;

b) procedure of escaping from a critical point, which is based on the global optimality conditions (GOC), with the following inclusion of the local search.

$$g(x) - h(x) \downarrow \min, \quad x \in D. \quad (\text{DC})$$

Theorem 3. (Global Optimality Conditions)

Let a feasible point $z \in D$ be a global solution to Problem (DC) ($z \in \text{Sol}(\text{DC})$). Then

$$\begin{aligned} (\mathcal{E}): \quad & \forall (y, \beta): \quad y \in D, \quad \beta - h(y) = \zeta \stackrel{\Delta}{=} g(z) - h(z), \\ & g(y) \leq \beta \leq \sup(g, D), \\ & g(x) - \beta \geq \langle h'(y), x - y \rangle, \quad \forall x \in D, \end{aligned} \quad (14)$$

where $h'(y) \in \partial h(y)$.

If, in addition, the following regularity condition holds

$$(\mathcal{H}): \quad \exists v \in D: \quad g(v) - h(v) > \zeta, \quad (15)$$

then conditions (\mathcal{E}) turns out to be sufficient for the point z being a global solution to Problem (DC).



Suppose, the 3-tuple $(\hat{y}, \hat{\beta}, \hat{x})$, such that $(\hat{y}, \hat{\beta}): h(\hat{y}) = \hat{\beta} - \zeta$, $\zeta := f(z)$ and $\hat{x} \in D$, violates the GOC (\mathcal{E}), i.e.

$$g(\hat{x}) < \hat{\beta} + \langle h'(\hat{y}), \hat{x} - \hat{y} \rangle,$$

Then from convexity of $h(\cdot)$ it follows that

$$f(\hat{x}) = g(\hat{x}) - h(\hat{x}) < h(\hat{y}) + \zeta - h(\hat{y}) = f(z)$$

or $f(\hat{x}) < f(z)$. Therefore, $\hat{x} \in D$ is „better“ than z .

And so, overhauling „perturbation parameters“ (y, β) in (\mathcal{E}) and solving linearized problems (see GOC)

$$g(x) - \langle h'(y), x \rangle \downarrow \min, \quad x \in D, \tag{16}$$

we obtain a family of initial points $x(y, \beta)$ for the local search methods.

$$g(x) - h(x) \downarrow \min, \quad x \in D. \quad (DC)$$

1) Find a critical point z with special local search algorithm, $\zeta \triangleq g(z) - h(z)$.

2) Choose number $\beta \in [\beta_-, \beta_+]$, $\beta_- = \inf(g, D)$, $\beta_+ = \sup(g, D)$.

3) Construct an approximation

$$A(\beta) = \{v^1, \dots, v^N \mid h(v^i) = \beta - \zeta, \quad i = 1, \dots, N, \quad N = N(\beta)\},$$

of level surface of function $h(\cdot)$.

4) Beginning from each point v^i of the approximation $A(\beta)$ find a point u^i by means of a special local search algorithm.

5) Verify the inequality from GOC:

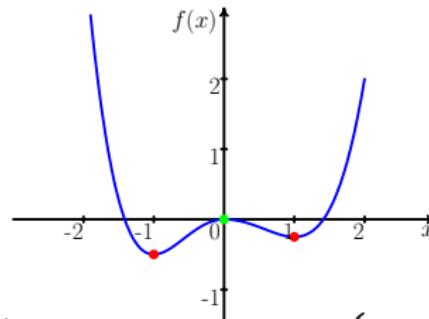
$$g(u^i) - \beta \geq \langle h'(v^i), u^i - v^i \rangle \quad \forall i = 1, 2, \dots, N.$$

$$(w^i \mapsto v^i)$$



Example 4

$$f(x) = \begin{cases} \frac{1}{4}x^4 - \frac{1}{2}x^2, & x \geq 0, \\ \frac{1}{2}x^4 - x^2, & x \leq 0. \end{cases} \quad (17)$$



$$g(x) = \begin{cases} g_1(x) = \frac{1}{4}x^4, & x \geq 0, \\ g_2(x) = \frac{1}{2}x^4, & x \leq 0; \end{cases} \quad h(x) = \begin{cases} h_1(x) = \frac{1}{2}x^2, & x \geq 0, \\ h_2(x) = x^2, & x \leq 0. \end{cases} \quad (18)$$

Global Optimality Conditions

$$\left. \begin{array}{l} (\mathcal{E}): \quad \forall (y, \beta): \quad y \in D, \quad \beta - h(y) = \zeta \triangleq g(z) - h(z), \\ \qquad \qquad \qquad g(y) \leq \beta \leq \sup(g, D), \\ (VI): \quad g(x) - \beta \geq \langle \nabla h(y), x - y \rangle, \quad \forall x \in D, \end{array} \right\}$$

How is efficient a local search method?

$s = 0$

$$x_0 = 100, \quad \nabla h(x_0) = x_0 = 100$$

$$(\mathcal{PL}_0): \quad \varphi_0(x) = g(x) - \langle \nabla h(x_0), x \rangle = \frac{1}{4}x^4 - 100x \downarrow \min_x, \quad x \in I\!\!R,$$

$$\nabla \varphi_0(x) = x^3 - 100 = 0, \quad 4^3 = 64, \quad 5^3 = 125, \quad x_1 \approx 4,5$$

$$(\mathcal{PL}_1): \quad \varphi_1(x) = \frac{1}{4}x^4 - 4,5x \downarrow \min_x, \quad x \in I\!\!R,$$

$s = 1$

$$\nabla \varphi_1(x) = x^3 - 4,5 = 0 \Rightarrow x_2 = 1,7,$$

$$(\mathcal{PL}_2): \quad \varphi_2(x) = \frac{1}{4}x^4 - 1,7x \downarrow \min_x, \quad x \in I\!\!R,$$

$s = 2$

$$\begin{aligned} \nabla \varphi_2(x) &= x^3 - 1,7 = 0 \Rightarrow x_3 = 1,2 \\ x_s &\rightarrow 1 \in Sol(\mathcal{P}). \end{aligned}$$

Example 4. Global Search Algorithm

Step 1. $x^0 = 100 \mapsto z^1 = 1, \zeta_1 = f(z^1) = -\frac{1}{4}$.

Step 2. $\beta_0 = g(z^k) = g(z^1) = g_1(1) = \frac{1}{4} \cdot 1 = \frac{1}{4}$.

Step 3. Construct the approximation

$$\mathcal{A}_k = \{y_1, y_2, \dots, y_N \mid h(y_i) = \beta_0 - \zeta_1 = \frac{1}{4} - \left(-\frac{1}{4}\right) = \frac{1}{2}\}.$$

i = 1. $h_1(y_1) \triangleq \frac{1}{2}y_1^2 = \frac{1}{2}, y_1^2 = 1, y_1 = 1$ (τ.κ. $y_1 \geq 0$).

Step 4.

$$g(x) - \langle \nabla h(y_1), x \rangle \downarrow \min_x, \quad x \in I\!\!R \quad (\mathcal{P}L_i)$$

$$\varphi_1(x) = \frac{1}{4}x^4 - \langle \nabla h(y_1), x \rangle = \frac{1}{4}x^4 - 1 \cdot x \downarrow \min_x,$$

$$\nabla \varphi_1(x) = x^3 - 1 = 0, \quad \bar{x}_1 = 1.$$

Step 5. (Verification of VI). $g(\bar{x}_1) - \beta_0 \stackrel{?}{\geq} \langle \nabla h(y_1), \bar{x}_1 - y_1 \rangle$,

$$\frac{1}{4}\bar{x}_1^4 - \frac{1}{4} = \frac{1}{4} - \frac{1}{4} = 0 \geq \langle 1, 1 - 1 \rangle = 0. VI \text{ is not violated! Go to step 3.}$$

i = 2. $h_2(y) = y^2 = \frac{1}{2}, y < 0, y_2 = -\frac{\sqrt{2}}{2}$.

Step 4.

$$\varphi_2(x) = g(x) - \langle \nabla h_2(y_2), x \rangle = \frac{1}{2}x^4 - \langle 2y_2, x \rangle = \frac{1}{2}x^4 - \sqrt{2} \cdot x \downarrow \min_x, \quad (\mathcal{P}L_2)$$

$$\nabla \varphi_2(x) = 2x^3 + \sqrt{2} = 0, \quad \bar{x}_2 = -\left(\frac{\sqrt{2}}{2}\right)^{\frac{1}{3}}.$$



Example 4. Global Search Algorithm

Step 5. (*Verification of VI*). $\nabla h_2(x) = 2x$

$$g_2(\bar{x}_2) - \beta_0 \geq \langle \nabla h_2(y_2), \bar{x}_2 - y_2 \rangle$$

$$\frac{1}{2}\bar{x}_2^4 - \frac{1}{4} \geq \langle 2y_2, \bar{x}_2 - y_2 \rangle$$

$$\frac{1}{2} \left(\frac{\sqrt{2}}{2} \right)^{\frac{4}{3}} - \frac{1}{4} \geq \langle -\sqrt{2}, -\left(\frac{\sqrt{2}}{2} \right)^{\frac{1}{3}} + \frac{\sqrt{2}}{2} \rangle = \langle \sqrt{2}, \left(\frac{\sqrt{2}}{2} \right)^{\frac{1}{3}} - \frac{\sqrt{2}}{2} \rangle > 0$$

$$\frac{1}{2} \left[\left(\frac{\sqrt{2}}{2} \right)^{\frac{4}{3}} - \frac{1}{2} \right] \geq \langle \sqrt{2}, \left(\frac{\sqrt{2}}{2} \right)^{\frac{1}{3}} - \frac{\sqrt{2}}{2} \rangle.$$

There is no violation of VI! Change β .

Step 2. $\beta = \frac{3}{4}$.

Step 3. Construct the approximation

$$h(y) = \begin{cases} \frac{1}{2}y^2, & y \geq 0, \\ y^2, & y < 0 \end{cases} = \beta_1 - \zeta_1 = \frac{3}{4} - \left(-\frac{1}{4} \right) = 1.$$

i = 1. $\frac{1}{2}y_1^2 = 1$, $y_1^2 = 2$, $y_1 = \sqrt{2}$.

i = 2. $y_2^2 = 1$, $y_2 = -1 < 0$,

Example 4. Global Search Algorithm

Step 4. (*Solving* $(\mathcal{P}L_i)$). $y_2 = -1 < 0$.

$$g(x) - \langle \nabla h(y_i), x \rangle = \varphi_i(x) \downarrow \min_x, \quad x \in I\!\!R \quad (\mathcal{P}L_i)$$

$i = 2$.

$$\varphi_2(x) = \frac{1}{2}x^4 - 2(-1)x = \frac{1}{2}x^4 + 2x \downarrow \min_x$$

$$\nabla \varphi_2(x) = 2x^3 + 2 = 2(x^3 + 1) = 0, \quad x^3 = -1, \quad \bar{x}_2 = -1.$$

Step 5. (*Verification of VI*). $\bar{x}_2 = -1$.

$$g_2(\bar{x}_2) - \beta_1 \geq \langle \nabla h_2(y_2), \bar{x}_2 - y_2 \rangle$$

$$\frac{1}{2}\bar{x}_2^4 - 1 \geq \langle \nabla h_2(y_2), -1 - (-1) \rangle = 0$$

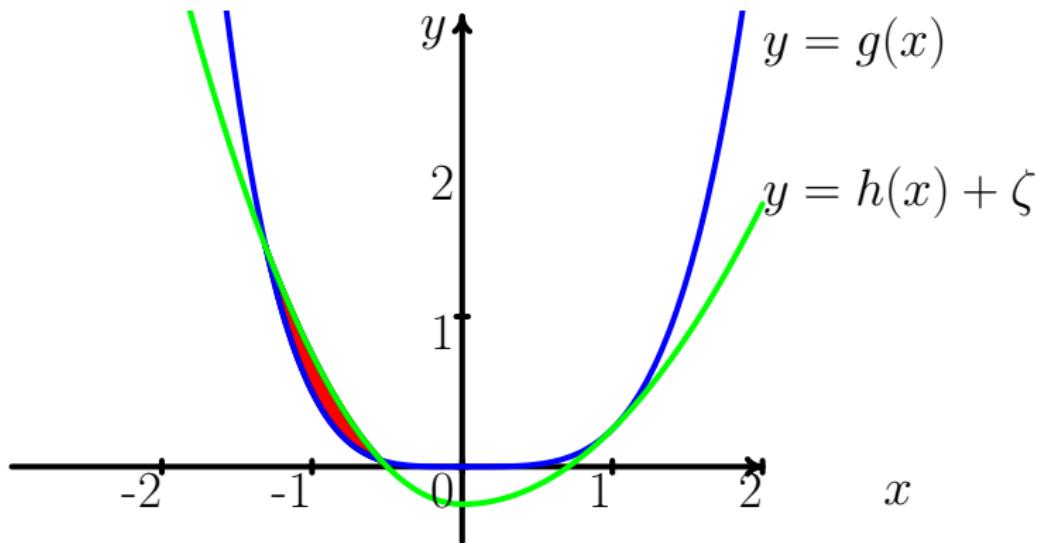
$$\frac{1}{2} - 1 = -\frac{1}{2} < 0.$$

VI is violated! Go to point $\bar{x}_2 = -1$,

$$\bar{x}_2 \in \operatorname{Arg\,min}(f, I\!\!R) = Sol(\mathcal{P}).$$

Example 4. Global Optimality Condition

$$\text{epi } g \subset \text{epi } (h(\cdot) + \zeta) \quad (\mathcal{E}')$$



Example 5

[http://www.mat.univie.ac.at/neum/glopt/coconut/Benchmark/
Library2_new_v1.html](http://www.mat.univie.ac.at/neum/glopt/coconut/Benchmark/Library2_new_v1.html)

“errinros-problem” ($n = 50$)

$$F(x) = \sum_{i=2}^n (x_{i-1} - 16\alpha_i^2 x_i^2)^2 + \sum_{i=2}^n (x_i - 1)^2 \downarrow \min_x.$$

$$F_{glob} = 39.904.$$

D.C. decomposition

$$F(x) = g(x) - h(x),$$

$$g(x) = \sum_{i=2}^n [x_{i-1}^2 + 256\alpha_i^2 x_i^2 + x_i^2 - 2x_i + 1 + 8(1+x_{i-1})^4 + 16x_i^2 + 8(x_i^2 + x_{i-1}^2)^2],$$

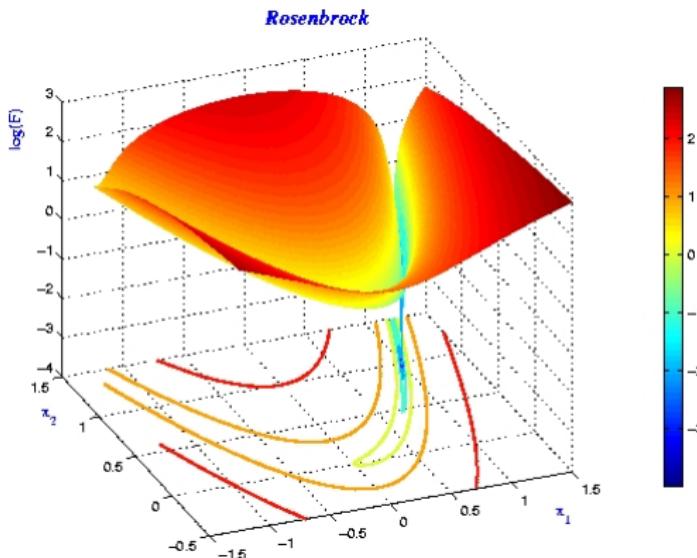
$$h(x) = \sum_{i=2}^n 8[(x_i^2 + (1+x_{i-1})^2)^2 + x_{i-1}^4].$$

$\ x_{glob} - x^0\ $	$F(x^0)$	St	PL	$\ x_{glob} - x_*\ $	$F(x_*)$	Time (cek)
1.032	40.504	2	4	$1.37 \cdot 10^{-10}$	39.904	4.51
3.174	67.943	8	32	$4.02 \cdot 10^{-5}$	39.907	27.05
17.904	370.531	42	398	$8.04 \cdot 10^{-9}$	39.906	63.92
59.32	$23 \cdot 10^9$	79	674	$9.012 \cdot 10^{-13}$	39.909	191.83
9.812	94.7	49	206	$2.04 \cdot 10^{-7}$	39.905	49.86
931.7	$12 \cdot 10^{25}$	107	2948	$3.07 \cdot 10^{-4}$	39.911	416.76
34.91	$77 \cdot 10^5$	12	82	$5.38 \cdot 10^{-7}$	39.906	82.97
107.23	$92 \cdot 10^{10}$	27	380	$4.28 \cdot 10^{-8}$	39.906	158.04
31.3	95982	18	138	$6.38 \cdot 10^{-8}$	39.905	73.53
97.21	$44 \cdot 10^3$	32	222	$2.94 \cdot 10^{-6}$	39.907	127.52
27.912	36654	12	72	$2.83 \cdot 10^{-8}$	39.906	44.75

$$F_{glob} = 39.904.$$

Minimization of Rosenbrock's function

$$(\mathcal{P}): \quad F(x) = \sum_{i=1}^{n-1} \left[(x_i - 1)^2 + N(x_{i+1} - x_i^2)^2 \right] \downarrow \min, \quad x \in I\!\!R^n.$$



$$F(x) = G(x) - H(x), \quad (19)$$

$$G(x) = \sum_{i=1}^{n-1} [(x_i - 1)^2 + \frac{N}{2}(1 + x_{i+1})^4 + \frac{3N}{2}x_i^4 + Nx_i^2x_{i+1}^2 + Nx_i^2 + Nx_{i+1}^2],$$

$$H(x) = \frac{N}{2} \sum_{i=1}^{n-1} ((1 + x_{i+1})^2 + x_i^2)^2.$$

$$G(x) = \sum_{i=1}^{n-1} g(x_i, x_{i+1}), \quad H(x) = \sum_{i=1}^{n-1} h(x_i, x_{i+1}). \quad (20)$$

The functions $G(\cdot)$ and $H(\cdot)$ are convex.

Local search testing for the problem of Rosenbrock's function minimization

n	Special Local Search Method				BFGS Method			
	$\ z - x_*\ $	$F(z)$	It	Time	$\ z - x_*\ $	$F(z)$	It	Time
5	1.12	0.64	46	1.25	2.01	3.93	75	0.04
5	1.12	0.64	44	1.25	2.01	3.93	65	0.04
5	1.12	0.64	44	1.35	2.01	3.93	90	0.06
5	$1.3 \cdot 10^{-3}$	$5.7 \cdot 10^{-5}$	56	1.30	$3.1 \cdot 10^{-5}$	$5.2 \cdot 10^{-9}$	61	0.02
10	1.22	0.81	45	1.30	1.99	3.98	67	0.55
10	1.22	0.81	44	1.30	1.99	3.98	75	0.55
10	$4.1 \cdot 10^{-4}$	$8.5 \cdot 10^{-6}$	49	1.40	1.99	3.98	71	0.50
10	0.024	0.0045	59	1.60	$1.4 \cdot 10^{-5}$	$3.33 \cdot 10^{-9}$	129	0.81
50	1.31	0.76	46	7.20	1.87	4.20	85	2.85
50	1.31	0.76	48	7.35	1.87	4.20	93	2.90
50	0.0035	$1.2 \cdot 10^{-5}$	48	7.40	1.87	4.20	88	2.85
50	$1.8 \cdot 10^{-4}$	$3.1 \cdot 10^{-6}$	57	8.25	$1.3 \cdot 10^{-6}$	$2.6 \cdot 10^{-10}$	136	3.15
100	1.46	4.56	51	14.20	1.73	0.83	112	10.00
100	1.46	4.56	50	13.90	1.73	0.83	120	10.30
100	0.0012	$4.5 \cdot 10^{-5}$	53	15.25	1.73	0.83	115	10.00
100	$3.8 \cdot 10^{-4}$	$5.5 \cdot 10^{-7}$	61	17.10	$1.8 \cdot 10^{-7}$	$8.2 \cdot 10^{-11}$	145	12.00

Global search testing for Rosenbrock's function minimization

n	$F(x^0)$	$\ \nabla F(x^0)\ $	$\ z - x_*\ $	$F(z)$	PL	$Time$
5	$1.2 \cdot 10^6$	408815	$1.1 \cdot 10^{-5}$	$1.6 \cdot 10^{-9}$	41	4.21
5	458810	196853	$1.2 \cdot 10^{-5}$	$7.5 \cdot 10^{-10}$	36	4.11
5	$2.5 \cdot 10^6$	617951	$1.2 \cdot 10^{-5}$	$1.7 \cdot 10^{-9}$	38	4.15
10	38084	24647	$3.2 \cdot 10^{-5}$	$2.0 \cdot 10^{-9}$	85	9.60
10	7940	8543	$3.1 \cdot 10^{-5}$	$1.8 \cdot 10^{-9}$	82	9.55
10	11101	11388	$1.5 \cdot 10^{-5}$	$1.8 \cdot 10^{-9}$	76	9.43
50	$8.1 \cdot 10^7$	$2.9 \cdot 10^6$	$3.4 \cdot 10^{-5}$	$2.5 \cdot 10^{-9}$	410	45.25
50	686176	256999	$1.7 \cdot 10^{-5}$	$3.7 \cdot 10^{-9}$	386	45.01
50	102509	32805.5	$2.1 \cdot 10^{-5}$	$4.2 \cdot 10^{-9}$	431	46.60
100	181996	44044.5	$1.6 \cdot 10^{-5}$	$2.5 \cdot 10^{-9}$	753	1 : 14.00
100	$2.1 \cdot 10^7$	$4.8 \cdot 10^6$	$2.2 \cdot 10^{-5}$	$2.6 \cdot 10^{-9}$	781	1 : 15.00
100	152183	41207	$2.4 \cdot 10^{-6}$	$6.8 \cdot 10^{-10}$	812	1 : 17.00
125	$4,4 \cdot 10^7$	$5,7 \cdot 10^6$	$3 \cdot 10^{-6}$	$1,5 \cdot 10^{-9}$	1044	1 : 50.10
125	$7,8 \cdot 10^7$	$4,1 \cdot 10^6$	$6,7 \cdot 10^{-6}$	$6,3 \cdot 10^{-9}$	1012	1 : 49.00
125	$3,6 \cdot 10^7$	$7,4 \cdot 10^6$	$4,6 \cdot 10^{-6}$	$2,1 \cdot 10^{-10}$	1101	1 : 51.00
151	$3.0 \cdot 10^8$	$7,2 \cdot 10^7$	$2,2 \cdot 10^{-5}$	$5,1 \cdot 10^{-9}$	1361	2 : 26.10
151	$5,8 \cdot 10^8$	$8,6 \cdot 10^7$	$3,8 \cdot 10^{-5}$	$3,1 \cdot 10^{-10}$	1355	2 : 25.20
151	$8,8 \cdot 10^8$	$8,1 \cdot 10^7$	$2,2 \cdot 10^{-5}$	$5,1 \cdot 10^{-9}$	1394	2 : 28.00

$$\|\nabla F(z)\| < 10^{-3}$$

For all dimensions the number of iteration of global search equals $It = 2$.

Problem Formulation

$$(DCC): \quad \begin{cases} f_0(x) \downarrow \min, & x \in S, \\ F(x) = g(x) - h(x) \leq 0. \end{cases}$$

Procedure 1 constructs $x(y) \in S$ from a predetermined point $y \in S$, $F(y) = g(y) - h(y) \leq 0$:

$$F(x(y)) = 0, \quad f_0(x(y)) \leq f_0(y).$$

Procedure 2 consists in sequential solving of the linearized problems

$$(LQ(u, \xi)): \quad g(x) - \langle \nabla h(u), x \rangle \downarrow \min, \quad x \in S, \quad f_0(x) \leq \xi. \quad (21)$$

$$(H_0): \quad \exists v \in S, \quad g(v) - h(v) > 0: \quad f_0(v) < f_0^* \stackrel{\triangle}{=} \mathcal{V}(P). \quad (22)$$

Theorem 4

Let $f_0(\cdot)$ and S be convex. Furthermore, $\mathcal{F}_0 = \{x \in S | f_0(x) \leq f_0(x_0)\}$ is bounded and (H_0) holds. Hence, the Special Local Search Method (SLSM) writes:

- i) after a finite number of iterations, we receive a point $y^N \in S$, $F(y^N) = 0$, which is ε_N -solution of problem $(LQ(y_N, \xi_N))$, where N is the number of the stop-iteration;
- ii) otherwise, for $\{x^s\}$ and $\{y^s\}$:

$$\begin{aligned} x^s &\in S, \quad F(x^s) = 0, \quad y^s \in S, \quad F(y^s) > 0, \\ \xi_{s+1} &:= f_0(x^{s+1}) < f_0(y^s) \leq \xi_s := f_0(x^s), \\ \xi_* &:= \lim_{s \rightarrow \infty} \xi_s = \lim_{s \rightarrow \infty} f_0(y^s), \end{aligned}$$

the following condition is satisfied: $0 = F(x^s) = \lim_{s \rightarrow \infty} F(y^s)$,

$x_* = \lim_{s \rightarrow \infty} x^s = \lim_{s \rightarrow \infty} y^s$, for some $x_* \in I\!\!R^n$, $F(x_*) = 0$.

Furthermore, x_* is a solution of the linearized problem $(LQ(x_*, \xi_*))$ and the normal stationary point for the dual problem $(Q(\xi_*))$.

D.C. constrained problem

$$\left. \begin{array}{l} f_0(x) \downarrow \min, \quad x \in S, \\ g(x) - h(x) \leq 0. \end{array} \right\} \quad (DCC)$$

Theorem 5. (Necessary GOC)

Let the following assumption (\mathcal{H}_1) be fulfilled:

$$\exists v \in S : F(v) > 0, \quad f_0(v) < \mathcal{V}(DCC) \triangleq \inf_x \{f_0(x) \mid x \in S, F(x) \leq 0\}.$$

If $z \in Sol(DCC)$, then ($h'(y) \in \partial h(y)$):

$$\left. \begin{array}{l} \forall (y, \beta) : h(y) = \beta, \quad y \in S, \\ g(x) - \beta \geq \langle h'(y), x - y \rangle \quad \forall x \in S : f_0(x) \leq f_0(z). \end{array} \right\} \quad (\mathcal{E}_1)$$

D.C. constrained problem

$$\left. \begin{array}{l} f_0(x) \downarrow \min, \quad x \in S, \\ g(x) - h(x) \leq 0. \end{array} \right\} \quad (DCC)$$

Theorem 6. (Sufficient GOC)

Let the following regularity conditions hold:

$$\exists v \in S : F(v) > 0, \quad (23)$$

$$\left. \begin{array}{l} \forall y \in S : F(y) = 0 \ (g(y) = h(y)) \\ \exists p = p(y) \in S : g(p) - g(y) < \langle h'(y), p - y \rangle. \end{array} \right\} \quad (\mathcal{H})$$

If $z \in S$, $F(z) = 0$ and, in addition, the condition

$$\left. \begin{array}{l} \forall(y, \beta) : h(y) = \beta, \ y \in S, \quad g(y) \leq \beta \leq \sup(g, S), \\ g(x) - \beta \geq \langle h'(y), x - y \rangle \quad \forall x \in S : f_0(x) \leq f_0(z), \end{array} \right\} \quad (\mathcal{E}_0)$$

is satisfied, then $z \in Sol(DCC)$.

Example 6

Consider the following problem ($x \in I\!\!R$)

$$f(x) = (x - 1)^3 \downarrow \min, \quad F(x) = |x| - \frac{x^2}{2} \geqslant 0. \quad (24)$$

Suppose $g(x) = |x|$, $h(x) = \frac{x^2}{2}$. It is easy to see $z = 0$ ($F(z) = 0$) is a critical point in the sense of the classical definition of optimality

$$\{\nabla f(z)\} \cap \{cone\{\partial g(z)\} - \{\nabla h(z)\}\} \neq \emptyset.$$

Let $\beta = 2$, $y = -2$, $g(y) = 2$, $\nabla g(y) = -1$. Let us consider also the point $u = -1.5$, that is a feasible point in the linearized problem

$$h(x) - \langle \nabla g(y), x \rangle \downarrow \min, \quad f(x) \leqslant f(z).$$

Then

$$h(u) - \beta - \langle \nabla g(y), u - y \rangle = \frac{1}{2}2.25 - 1.5 < 0.$$

Therefore, the point z is not global solution of problem (24).

Now we prove that $z = -2$ is a global solution.

$$f(z) = -27 \stackrel{\triangle}{=} \rho, \quad F(z) = 0.$$

Example 6

Continuation

Consider linearized problem for all pair (y, β) such that $g(y) \triangleq |y| = \beta$,
 $y^* \in \partial g(y)$:

$$h(x) - \langle y^*, x \rangle \downarrow \min, \quad f(x) \leq \rho = -27. \quad (25)$$

We recall that

$$y^* = \begin{cases} 1, & \text{if } y > 0 \quad (y - \beta = 0), \\ -1, & \text{if } y < 0 \quad (y + \beta = 0), \\ \alpha \in [-1, 1], & \text{if } y = 0 \quad (\beta = 0). \end{cases}$$

It is easy to see that the point $\hat{x} = -2$ is a unique solution of problem (25) for all $y^* \in [-1, 1]$.

Further we verify the condition (\mathcal{E}_0) in all cases.

a) $y^* = 1$ ($y > 0, y - \beta = 0$). By inequality $\beta \geq h(y)$, we consider only $y \leq 2$.

$$h(\hat{x}) - \beta - \langle y^*, \hat{x} - y \rangle = 4 - \beta + y = 4 > 0.$$

b) $y^* = -1$ ($y < 0, y + \beta = 0$),

$$h(\hat{x}) - \beta - \langle y^*, \hat{x} - y \rangle = 2 - \beta - 2 - y = 0.$$

c) $y^* = 0$ ($y = 0 = \beta$),

$$h(\hat{x}) - \beta - \langle y^*, \hat{x} - y \rangle = 2 > 0.$$

d) $y^* = \alpha \in]0, 1[$ ($y = 0 = \beta$),

$$h(\hat{x}) - \beta - \langle y^*, \hat{x} - y \rangle = 2 - \alpha(-2) > 0.$$

e) $y^* = -\alpha, \alpha \in]0, 1[$ ($y = 0 = \beta$),

$$h(\hat{x}) - \beta - \langle y^*, \hat{x} - y \rangle = 2(1 + \alpha) > 0.$$

Example 6

Continuation

Therefore, the basic inequality (\mathcal{E}_0) holds in all cases. Nevertheless, it is not clear that $z = -2$ is global solution, because not all of the assumptions of the Theorem 5 are verified. It is easy to prove that (12) holds, for instance, with $v = 3$. To verify the hypothesis (\mathcal{H}) , we consider the problem without restrictions:

$$h(x) - \langle y^*, x \rangle \downarrow \min, \quad x \in I\!\!R, \quad (26)$$

$y^* \in \partial g(y)$ for all y such that $F(y) = 0$, i.e. $\frac{y^2}{2} = |y|$. Hence, it suffices to verify the inequality $h(x(y)) - h(y) - \langle y^*, x(y) - y \rangle < 0$ with $y_1 = 2$, $y_2 = -2$, $y_3 = 0$, where $x(y)$ is a solution of (26), corresponding y_i , and $y^* \in \partial g(y_i)$, $i = 1, 2, 3$. It can be proved that $x(y) = y^*$.

If the last inequality holds, then (\mathcal{H}) is fulfilled with $p(y) = x(y)$.

We get three cases:

$$1) \quad y_i = 2, \quad y_i^* = 1 = x_1 \stackrel{\Delta}{=} x(y_1), \quad h(x_1) - h(y_1) - \langle y_1^*, x_1 - y_1 \rangle = \frac{-1}{2} < 0;$$

$$2) \quad y_2 = -2, \quad y_2^* = -1 = x_2 \stackrel{\Delta}{=} x(y_2), \quad h(x_2) - h(y_2) - \langle y_2^*, x_2 - y_2 \rangle = \frac{-1}{2} < 0;$$

$$3) \quad y_3 = 0, \quad y^* \in \partial g(y_3) = [-1, 1]. \quad x_3 \stackrel{\Delta}{=} x(y_3) = y_3^* = \alpha \in [-1, 1], \text{ and if } \alpha \neq 0, \\ \text{then } h(x_3) - h(y_3) - \langle y_3^*, x_3 - y_3 \rangle = \frac{-\alpha^2}{2} < 0.$$

Therefore, the inequality in (\mathcal{H}) is fulfilled $\forall \alpha \neq 0$, while it is sufficient that this inequality for one α . The assumptions (\mathcal{H}) and (12) hold. Thus $z = -2$ is a global solution of (24).

Example 7

$$\begin{cases} f(x) = \frac{1}{2}(x_1 - 4)^2 + (x_2 + 2)^2 \downarrow \min, \\ F(x) = (x_2 + 1)^2 - (x_1 - 1)^2 \geq 0. \end{cases} \quad (27)$$

It is easy to see that the point $z = \left(\frac{4}{3}, \frac{-2}{3}\right)$ ($F(z) = 0$) is stationary, i.e.

$$\nabla f(z) - \lambda \nabla F(z) = 0, \quad \lambda F(z) = 0,$$

при $\lambda = 4$. Let $g(x) = (x_2 + 1)^2$, $h(x) = (x_1 - 1)^2$.

Further, we prove that z is not a global solution of (27). Applying the condition (E)-(??) with $\beta = 4$, $y = (\alpha, -3)$, where $\alpha \in I\!\!R$, we receive

$$g(y) = \beta, \quad \nabla g(y) = (0, -4), \quad h(y) = (\alpha - 1)^2 \leq 4.$$

Moreover, consider the point

$$u = \left(4, -2 - \frac{4\sqrt{3}}{3}\right), \quad f(u) = \frac{16}{3} = f(z).$$

$$\text{Then } h(u) - \beta - \langle \nabla g(y), u - y \rangle = 5 + 4 \cdot \left(1 - \frac{4\sqrt{3}}{3}\right) < 0.$$

Therefore, according to the Theorem 4, the point z is not a global solution of (27). For global search it is necessary to implement a local search from the

feasible point $u = \left(4, -2 - \frac{4\sqrt{3}}{3}\right)$, such that $F(u) = (-3.3094)^2 - 9 > 0$.

THANK YOU!



THANK YOU!

